

Rate control for communication networks

shadow prices, proportional fairness
and stability

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Introduction



- ◆ This paper analyses the stability and fairness of two classes of rate control algorithms for communication networks.
- ◆ The algorithms provide natural generalizations to large-scale networks of simple additive increase/multiplicative decrease schemes

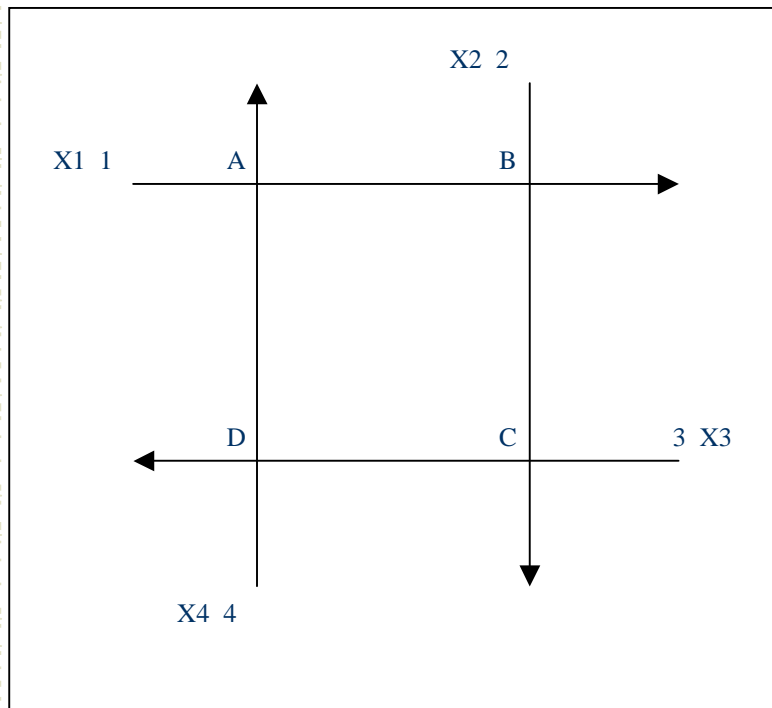
Model

Definition:

Consider a network with a set J of resources (abstraction of links), and let C_j be the finite capacity of resources j . Let a route r be a non-empty subset of J , and write R for the set of possible routes. Set $A_{jr}=1$ if $j \in r$ and $A_{jr}=0$ otherwise. This defines a 0-1 matrix A .

Associate a route r with a user, we allocate a rate x_r and a utility $U_r(x_r)$. Assuming $U_r(x_r)$ is increasing and strictly concave.

Model



- ◆ $J = \{A, B, C, D\}$
- ◆ $R = \{1, 2, 3, 4\}$
- ◆ $1 = \{A, B\}$

$$A = \begin{pmatrix} 1 & 0 & 0 & 1 \\ 1 & 1 & 0 & 0 \\ 0 & 1 & 1 & 0 \\ 0 & 0 & 1 & 1 \end{pmatrix} \begin{matrix} A \\ B \\ C \\ D \end{matrix}$$

1 2 3 4

- ◆ Users' rate: x_1, x_2, x_3, x_4

System Model

- ◆ SYSTEM(U, A, C)

$$\max \sum_{r \in R} U_r(x_r)$$

Subject to

$$Ax \leq C$$

over

$$x \geq 0$$

System Model

- ◆ This optimization problem is mathematically fairly tractable (with a strictly concave objective function and a convex feasible region), but it involves utilities U that are unlikely to be known by the network. We are thus led to consider two simpler problems.

User and Network Model

- ◆ Suppose that user r may choose an amount to pay per unit time, w_r , and receives in return a flow x_r proportional to w_r , say $x_r = w_r / \lambda_r$. We can get user model and network model:

User and Network Model

*USER*_{*r*} ($U_r; \lambda_r$):

$$\max U_r \left(\frac{w_r}{\lambda_r} \right) - w_r$$

over

$$w_r \geq 0$$

NETWORK ($A, C; w$):

$$\max \sum_{r \in R} w_r \log x_r$$

subject

$$Ax \leq C$$

over

$$x \geq 0$$

User and Network Model

- ◆ For User Model, the purpose is to maximize the user's gain (utility-payment).
- ◆ For Network, the purpose we choose Network Model as: $\max \sum_{r \in R} w_r \log x_r$ is to achieve fair. Define that the rates per unit charge of network is proportionally fair if x is feasible and for any other x^*

$$\sum_{r \in R} w_r \frac{x_r^* - x_r}{x_r} \leq 0 \Rightarrow \sum_{r \in R} w_r \frac{\delta x_r}{x_r} \leq 0$$

User and Network Model

- ◆ It means the smaller flows and more generous users will be favored.

- ◆ If x is a optimum solution to NETWORK, we have

$$(x_r, r \in R) \rightarrow (x_r + \delta x_r, r \in R) \Rightarrow \sum_{r \in R} w_r \frac{\delta x_r}{x_r} \leq 0$$

It is same as the above condition.

So when we maximize NETWORK, it is also insure that the rates per unit charge of network is proportionally fair.

User and Network Model

- ◆ It is proved in Kelly's paper "Charging and rate control for elastic traffic" that there always $\exists \lambda, w, x$, satisfying $w_r = \lambda_r x_r$ such that w_r solves USER_r and vector x solves NETWORK. Further x is the unique solution to SYSTEM and rates per unit charge of the NETWORK is proportionally fair.

User and Network Model

- ◆ Under the decomposition of the problem SYSTEM into the problems NETWORK and USER_r, the utility function U_r is not required by the network.

So suppose network knows w_r , we design algorithms to solve the NETWORK problem.

Primal Algorithm

$$\frac{d}{dt} x_r(t) = k \left(w_r - x_r(t) \sum_{j \in r} \mu_j(t) \right) \quad (1)$$

Where

$$\mu_j = p_j \left(\sum_{s=j \in s} x_s(t) \right) \quad (2)$$

$p_j(y)$ means that when the total flow of j is y ,
the unit flow charge is p_j

Primal Algorithm

- ◆ Define
$$u(x) = \sum_{r \in R} w_r \log x_r - \sum_{j \in J} \int_0^{x_j} p_j(y) dy$$
- ◆ Theorem 1: The strictly concave function $u(x)$ is a Lyapunov function for the system of differential equations (1)-(2). The unique value x maximizing $u(x)$ is a stable point of the system, to which all trajectories converge.

Primal Algorithm

- ◆ Proof: $u(x)$ is strictly concave on $x \geq 0$ with an interior maximum.

$$\frac{\partial}{\partial x_r} u(x) = \frac{\omega_r}{x_r} - \sum_{j \in r} p_j \left(\sum_{s: j \in s} x_s \right)$$

$$\frac{d}{dt} u(x(t)) = k \sum_{r \in R} \frac{1}{x_r(t)} \left(w_r - x_r(t) \sum_{j \in r} p_j \left(\sum_{s: s \in j} x_s(t) \right) \right)^2$$

establishing that $u(x)$ is strictly increasing with t , unless the unique x maximizing $u(x)$. The function $u(x)$ is thus a Lyapunov function for system above.

Primal Algorithm

- ◆ Define the continuous function:

$$p_j(y) = \frac{(y - C_j + \varepsilon)^+}{\varepsilon^2} = \mu_j$$

$\varepsilon \rightarrow 0 \Rightarrow$ maximization of the lyapunov function approximates arbitrarily closely the primal problem NETWORK.

Rate of convergence

◆ Let $x_r(t) = x_r + x_r^{\frac{1}{2}} y_r(t)$, linearizing

$$\frac{d}{dt} y(t) = -k \left(W X^{-1} + X^{\frac{1}{2}} A^T P' A X^{\frac{1}{2}} \right) y(t)$$

$$X = \text{diag}(x_r, r \in R) \quad W = \text{diag}(w_r, r \in R) \quad P' = \text{diag}(P'_j, j \in J)$$

$$\Gamma^T \Phi \Gamma = W X^{-1} + X^{\frac{1}{2}} A^T P' A X^{\frac{1}{2}}$$

$$\Gamma^T \Gamma = I \quad \Phi = \text{diag}(\Phi_r, r \in R)$$

we get
$$\frac{d}{dt} y(t) = -K \Gamma^T \Phi \Gamma y(t)$$

Rate of convergence

- ◆ Rate of convergence is determined by the smallest eigenvalue, Φ_r . And speed of convergence increases both with gain parameter K and P' .

| | |
|---------------|---------------------------------|
| $K \uparrow$ | Speed of convergence \uparrow |
| $P' \uparrow$ | Speed of convergence \uparrow |

Table 1

Stochastic analysis

- ◆ We consider a stochastic perturbation of linearized equation. Let

$$dy(t) = -\kappa(\Gamma^T \Phi \Gamma y(t) dt + F dB(t))$$

$B(t)$ is a collection of independent standard Brownian motions.

The stationary solution is:

$$y(t) = -\kappa \int_0^\infty e^{-\kappa(t-\tau)\Gamma^T \Phi \Gamma} F dB(\tau)$$

It has a multivariate normal distribution

$$y(t) \sim N(0, \Sigma)$$

Stochastic analysis

- ◆ So if we define

$$[\Gamma F; \Phi]_{rs} = \left[\int_{-\infty}^0 e^{\tau\Phi} \Phi F F^T \Gamma^T e^{\tau\Phi} d\tau \right]_{rs} = \frac{[\Gamma F F^T \Gamma^T]_{rs}}{\Phi_r + \Phi_s}$$

We get

$$\Sigma = k\Gamma^T [\Gamma F; \Phi] \Gamma$$

| | |
|---------------|---------------------|
| $K \uparrow$ | $\Sigma \uparrow$ |
| $P' \uparrow$ | $\Sigma \downarrow$ |

Table 2

As P' increases, not only is convergence faster, but also the spread at equilibrium decreases.

Various stochastic Model

- ◆ We have various sources of randomness that may lead to different covariance structure.
- ◆ Congestion indication with join feedback:

$$dx_r(t) = k \left(\omega_r dt - x_r(t) \sum_{j \in r} \varepsilon_j dN_j \left(\varepsilon_j^{-1} \int_0^t \mu_j(\tau) d\tau \right) \right)$$

Various stochastic Model

- ◆ Resource j generates feedback signal indicating congestion as a time-dependent Poisson process at rate $\epsilon_j^{-1} \mu_j(t)$
suppose that when a feedback signal is generated, it is send to each user r whose route passes through resource j , and the rate x_r reduces $K \epsilon x_r(t)$

Various stochastic Model

- ◆ We can get the corresponding Brownian version of the linearized system let

$$F_{rj} = \varepsilon_j^{\frac{1}{2}} \mu_j^{\frac{1}{2}} A_{jr} x_r^{\frac{1}{2}}$$

$$FF^T = X^{\frac{1}{2}} A^T E P A X^{\frac{1}{2}}$$

Various stochastic Model

- ◆ Congestion indication with individual feedback:

$$dx_r(t) = k \left(\omega_r dt - \sum_{j \in r} \varepsilon_j dN_{jr} \left(\varepsilon_j^{-1} \int_0^t x_r(\tau) \mu_j(\tau) d\tau \right) \right)$$

Feedback signals from resource j to user r arise at rate $\varepsilon_j^{-1} x_r(t) \mu_j(t)$. And user r reacts to such a feedback signal by reducing x_r with $k\varepsilon$.

We get

$$F_{r,(j,s)} = \varepsilon_j^{\frac{1}{2}} \mu_j^{\frac{1}{2}} A_{jr} \mathbf{I}[r = s]$$

Time Lags

Consider next the lagged, discrete time system

$$x_r[t+1] = x_r[t] + k \left(w_r - x_r[t] \sum_{j \in r} \mu_j[t - d(j, r)] \right)$$

$$\mu_j[t] = p_j \left(\sum_{s: j \in s} x_s[t - d(j, s)] \right)$$

$d(j, s)$ is the delay between user s changing its rate and the altered flow reaching resource j .

$d(j, r)$ is the delay between resource j generating a feedback signal and user r receiving it

Discrete System with Time Lags

Theorem 2:

The vector x maximizing the strictly concave function $u(x)$ is the unique equilibrium point of this system.

- ◆ Proof: The vector x is an equilibrium point if and only if it solves

$$w_r = x_r \sum_{j \in r} p_j \left(\sum_{s: j \in s} x_s \right)$$

It is precisely the stationarity condition of system (1)-(2)

Discrete System with Time Lags

- ◆ If we let $\mu_j = p_j \left(\sum_{s:j \in s} x_s \right)$, and suppose p_j is differentiable at point x_r . Let $x_r[t] = x_r + x_r^{1/2} y_r[t]$. Then, linearizing the above system about x .

If we define

$$(L[d])_{rs} = \sum_j p_j' A_{jr} A_{js} x_r^{1/2} x_s^{1/2} \mathbf{I}[d(j,r) + d(j,s) = d]$$

$$\sum_{d=0}^D L[d] = X^{1/2} A^T P' A X^{1/2}$$

Discrete System with Time Lags

$$L = \begin{pmatrix} \mathbf{I} - k(WX^{-1} + L[0]) & -kL[1] & -kL[2] & \dots & -kL[D] \\ \mathbf{I} & 0 & 0 & \dots & 0 \\ 0 & \mathbf{I} & 0 & \dots & \dots \\ \dots & \dots & \dots & \dots & \dots \\ 0 & 0 & 0 & \dots & 0 \end{pmatrix}$$

$$D = \max_{j,r,s} \{d(j,r) + d(j,s)\}$$

Discrete System with Time Lags

$$\begin{pmatrix} y[t+1] \\ y[t] \\ \dots \\ y[t-D+1] \end{pmatrix} = L \begin{pmatrix} y[t] \\ y[t-1] \\ \dots \\ y[t-D] \end{pmatrix}$$

- ◆ The equilibrium point x of the system is stable if and only if the spectral radius of the matrix L is less than unity.

Discrete System with Time Lags

| | |
|---------------|--------------------------|
| $K \uparrow$ | The system goes unstable |
| $P' \uparrow$ | The system goes unstable |

Table 3

Dual System

- ◆ The Lagrangian for the problem NETWORK is

$$L(x, z; \mu) = \sum_{r \in R} w_r \log x_r + \mu^T (C - Ax - z)$$

where $z \geq 0$ is a vector of slack variable and μ

is a vector of lagrange multipliers (or shadow price). Then

$$\frac{\partial L}{\partial x_r} = \frac{w_r}{x_r} - \sum_{j \in J} A_{jr} \mu_{jr}$$

So the unique optimum to the primal problem is given by

$$x_r = \frac{w_r}{\sum_{j \in J} A_{jr} \mu_{jr}}$$

Dual System

Where x_r, μ_j solve

$$\mu \geq 0, Ax \leq C, \mu^T (C - AX) = 0$$

Maximizing the Lagrangian over x and z gives

$$\begin{aligned} L(x, z; \mu) &= \sum_{r \in R} w_r \log \left(\frac{w_r}{\sum_{j \in J} A_{jr} \mu_{jr}} \right) + \mu^T C - \mu^T Ax \\ &= \sum_{r \in R} w_r \log(w_r) - \sum_{r \in R} w_r \log \left(\sum_{j \in J} A_{jr} \mu_{jr} \right) + \sum_{j \in J} \mu_j C_j - \sum_{r \in R} w_r \end{aligned}$$

Dual System

From it, removing the terms not dependent on the shadow price μ_j we can derive the DUAL:

$$\max \sum_{r \in R} w_r \log \left(\sum_{j \in J} \mu_j \right) - \sum_{j \in J} \mu_j C_j$$

Dual System

- ◆ Also we have a dual algorithm is:

$$\frac{d}{dt} \mu_j(t) = k \left(\sum_{r: j \in r} x_r(t) - q_j(\mu_j(t)) \right) \quad \underline{(3)}$$

$$x_r(t) = \frac{w_r}{\sum_{k \in r} \mu_k(t)} \quad (4)$$

$q_j(\eta)$ means the flow through resource j is $q_j(\eta)$, it generates a price η at resource j

Dual System

- ◆ If we define

$$v(\mu) = \sum_{r \in R} w_r \log \left(\sum_{j \in r} \mu_j \right) - \sum_{j \in J} \int_0^{\mu_j} q_j(\eta) d\eta$$

- ◆ Theorem 3: The strictly concave function v is a Lyapunov function for the system. The unique value μ maximizing v is a stable point of the system, to which all trajectories converge.

Dual System

- ◆ We also have rate of convergence:

$$\frac{d}{dt} \xi(t) = -k \Theta^T \Psi \Theta \xi(t)$$

| | |
|---------------|---------------------------------|
| $K \uparrow$ | Speed of convergence \uparrow |
| $Q' \uparrow$ | Speed of convergence \uparrow |

Table 4

Dual System

- ◆ Stochastic analysis: $d\xi(t) = -k(\Theta^T \Psi \Theta \xi(t) dt - G dB(t))$

let
$$[\Theta G; \Psi]_{jk} = \frac{[\Theta G G^T \Theta^T]_{jk}}{\psi_j + \psi_k}$$

then the covariance matrix is $\Sigma = k \Theta^T [\Theta G; \Psi] \Theta$

| | |
|---------------|---------------------|
| $K \uparrow$ | $\Sigma \uparrow$ |
| $Q' \uparrow$ | $\Sigma \downarrow$ |

Table 5

Dual System

- ◆ Discrete system with time lags:

$$\mu_j[t+1] = \mu_j[t] + k \left(\sum_{r:j \in r} x_r[t - d(j, r)] - q_j(\mu_j[t]) \right)$$

$$x_r[t] = \frac{w_r}{\sum_{k \in r} \mu_k[t - d(k, r)]}$$

Dual System

$$\begin{pmatrix} \xi[t+1] \\ \xi[t] \\ \dots \\ \xi[t-D+1] \end{pmatrix} = M \begin{pmatrix} \xi[t] \\ \xi[t-1] \\ \dots \\ \xi[t-D] \end{pmatrix}$$

$$(M[d])_{jk} = \sum_r x_r A_{jr} A_{kr} \mathbf{I}[d(j,r) + d(k,r) = d]$$

Dual System

$$M = \begin{pmatrix} I - k(M[0] + Q') & -kM[1] & -kM[2] & \dots & -kM[D] \\ I & 0 & 0 & \dots & 0 \\ 0 & I & 0 & \dots & 0 \\ \dots & \dots & \dots & \dots & \dots \\ 0 & 0 & 0 & \dots & 0 \end{pmatrix}$$

| | |
|---------------|--------------------------|
| $K \uparrow$ | The system goes unstable |
| $Q' \uparrow$ | The system goes unstable |

Table 6

Four nodes network

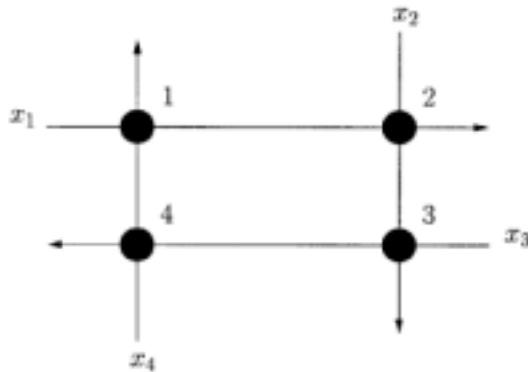


Figure 1 A four node network.

Suppose that the total load y on a resource takes the form of a Poisson stream of cells at rate y/ε .

Suppose that the time-axis is divided into non-overlapping slot $\tau\varepsilon$, and

that a feed back signal is generated for a slot if the total number of cells arriving in that slot exceeds a threshold N .

Four nodes network

- ◆ Joint feedback: suppose that when a feedback signal is generated, it is sent to each user r whose route passes through resource j , and the rate x_r reduces $\kappa \mathcal{E} x_r(t)$
- ◆ Individual feedback Suppose that when a feedback signal is generated at resource j , it is directed at a random route r with probability $\frac{x_r}{y}$, and rate x_r reduced $\kappa \mathcal{E}$

Four nodes network

- ◆ Consider the network illustrated above,
where $|J| = |R| = 4$
let $w_r = 0.0002$,and choose $N = 128$
 $\tau = 50$ so that $x_r = 1.0$ and $\mu_j = 0.0001$

Four nodes network

$$\Sigma_1 = \kappa \varepsilon 10^{-2} \begin{pmatrix} 2.4 & 0.8 & -0.8 & 0.8 \\ 0.8 & 2.4 & 0.8 & -0.8 \\ -0.8 & 0.8 & 2.4 & 0.8 \\ 0.8 & -0.8 & 0.8 & 2.4 \end{pmatrix}$$

$$\Sigma_2 = \kappa \varepsilon 10^{-1} \begin{pmatrix} 1.5 & -1.2 & 1.1 & -1.2 \\ -1.2 & 1.5 & -1.2 & 1.1 \\ 1.1 & -1.2 & 1.5 & -1.2 \\ -1.2 & 1.1 & -1.2 & 1.5 \end{pmatrix}$$

Four nodes network

- ◆ Compare the magnitudes and structures of the two matrices.
- ◆ We can see Σ_2 is larger in magnitude, since with individual feedback there are additional sources of variation in the random choice of which rate is to be reduced.

Four nodes network

- ◆ Note also that rates on routes sharing a node are positively correlated in joint feed back, because congestion indication at a node causes both routes through that node to decrease. However, for individual feedback, routes sharing a node are negatively correlated. A decrease in the flow on a route will allow increase on routes sharing a common node with it.

A random network

- ◆ We consider a network where the elements of the matrix A are independent random variables, each taking the value 1 with probability p and the value 0 otherwise.
- ◆ Let $w_r = \sum_j A(j, r)$ and $q_j(\eta) = \eta \sum_s A(j, s)$ so the unique stable point for the system (3)-(4) is $x_r = 1$ and $\mu_r = 1$

A random network

- ◆ Consider the system:

$$\mu_j[t+1] = \mu_j[t] + \kappa \left(\sum_{r: j \in r} N_r[t - d(j, r)] - q_j(\mu_j[t]) \right)$$

- ◆ Where $N_r[t]$ are collection of independent Poisson random variables, and has mean $x_r[t]$. Choose five random routes and resources for the following parameter: $J=100$, $R=1000$, $p=0.1$, $k=0.005$.

A random network

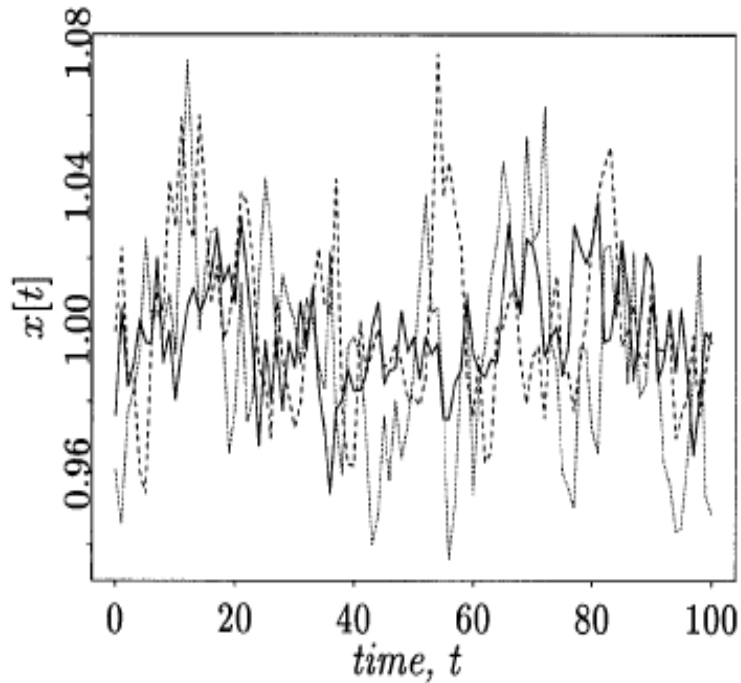


Figure 2 Rates on three randomly chosen routes.

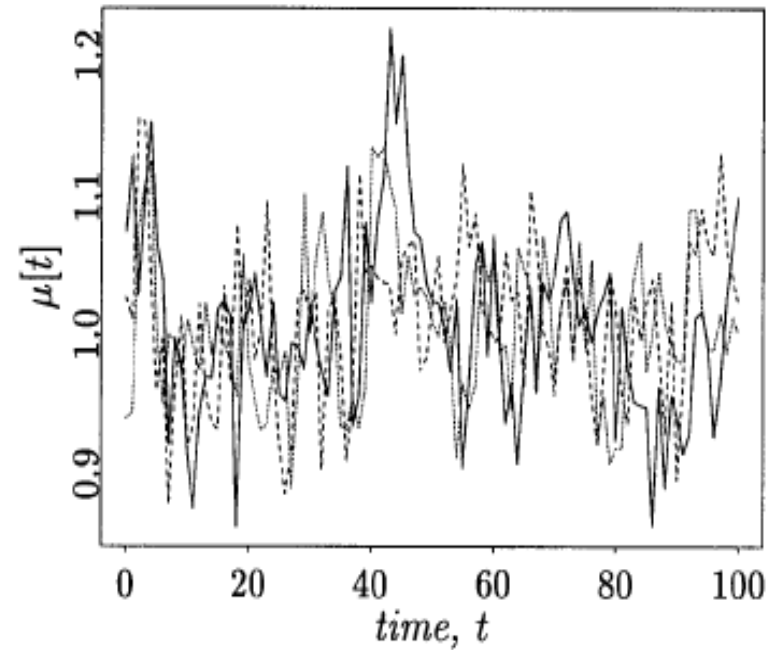


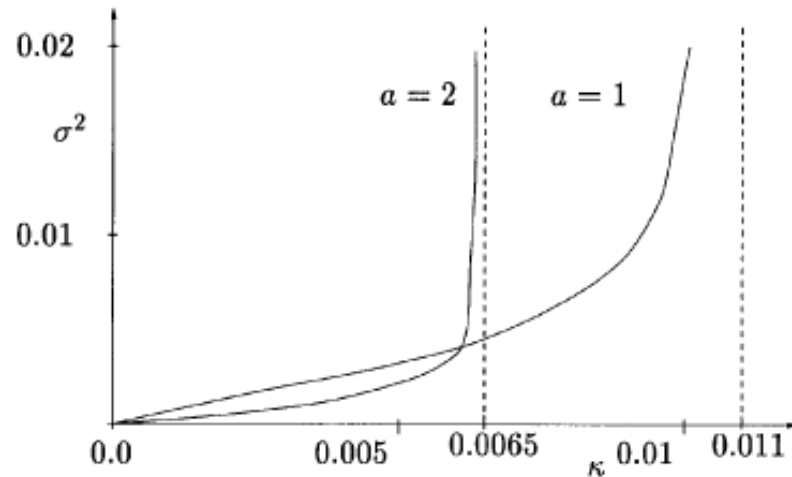
Figure 3 Shadow prices for three randomly chosen resources.

A random network

- ◆ From the above figures, we can see that for this example rates oscillate within a narrower band than shadow prices, and both are relatively well controlled.
- ◆ Finally, let us consider briefly the effect of more general choices for q_j , let

$$q_j = (a(\eta - 1) + 1) \sum_s A(j, s)$$

A random network



The case $a=1$ is discussed so far; the case $a=2$, corresponds to a doubling of the matrix of derivatives Q' . As predicted, increasing a is

to reduce variability, but also to lower the critical value of k at which the system becomes unstable

User adaptation

- ◆ Suppose user r is able to monitor its rate $x_r(t)$ and vary $w_r(t)$ to track the optimum of USER r . so we will always have $w_r = x_r U_r'(x_r)$

For the primal algorithm:

$$\frac{d}{dt} x_r(t) = k \left(w_r(t) - x_r(t) \sum_{j \in r} w_j(t) \right)$$

User adaptation

- ◆ The revised differential equation

$$u(x) = \sum_{r \in R} U_r(x_r) - \sum_{j \in J} \int_0^{x_j} p_j(y) dy$$

- ◆ Provides a Lyapunov function for the revised system, and the unique value maximizing $u(x)$ is a stable point of the system, to which all trajectories converge.

User adaptation

- ◆ Linearization may again be used to investigate behavior near the stable point:

$$\frac{d}{dt} y(t) = -kX^{\frac{1}{2}} \left(A^T P' A - U'' \right) X^{\frac{1}{2}} y(t)$$

$$U'' = \text{diag} \left(U_r''(x_r), r \in R \right)$$



End

Thanks!